

# Artificial Intelligence

**CSC 665**

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# **Search III**

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# Recap

- **Search:** make decisions by looking ahead
- **Logic:** deduce new facts from existing facts
- **Constraints:** find a way to satisfy a given specification
- **Probability:** reason quantitatively about uncertainty
- **Learning:** make future predictions from past observations

# Search

**Modeling:** start state, actions, cost, transition model, goal test

**Inference:** backtracking, DFS, BFS, UCS — *all uninformed search algorithms*

# Backtracking search (last time)

**Global state:** minimum cost path, set of explored nodes

**function** search( $s$ , path) :

- **if** IsEnd( $s$ ) :
  - update the minimum cost path
- **for each** action  $a \in \text{Actions}(s)$  :
  - **if** Succ( $s, a$ ) hasn't been explored yet:
    - add it to the explored set
    - extend path with Succ( $s, a$ ) and Cost( $s, a$ )
    - recurse: search(Succ( $s, a$ ), path)

# Backtracking search (last time)

**Global state:** minimum cost path, set of explored **(node, cost)** pairs

**function** search( $s$ , path) :

- **if** IsEnd( $s$ ) :
  - update the minimum cost path
- **for each** action  $a \in \text{Actions}(s)$  :
  - **if** Succ( $s, a$ ) hasn't been explored at **Cost( $s, a$ )** yet:
    - add **(Succ( $s, a$ ), Cost( $s, a$ ))** to the explored set
    - extend path with Succ( $s, a$ ) and Cost( $s, a$ )
    - recurse: search(Succ( $s, a$ ), path)

# Uniform Cost Search (UCS, Dijkstra's Algorithm)

- Start with a frontier that contains  $s_0$ , and an empty set of explored nodes
- **While** the frontier is nonempty:
  - Pop the node  $s$  with smallest priority  $p$  from the frontier
  - **If** IsEnd( $s$ ) : **return** solution
  - Add  $s$  to the explored set
  - **For** each  $a \in \text{Actions}(s)$ ,
    - Get  $s' = \text{Succ}(s, a)$
    - **If**  $s'$  is already explored: **continue**
    - Add  $s'$  to frontier with priority  $p + \text{Cost}(s, a)$



***[UCS example on board]***

# Correctness of UCS

**Theorem:** Assume action costs are non-negative. If a node  $s$  is popped from the frontier with priority  $p$ , then  $p$  is the cost of the min-cost path from  $s_0$  to  $s$ .

**Proof:** Take CSC 510 (or come to office hours).

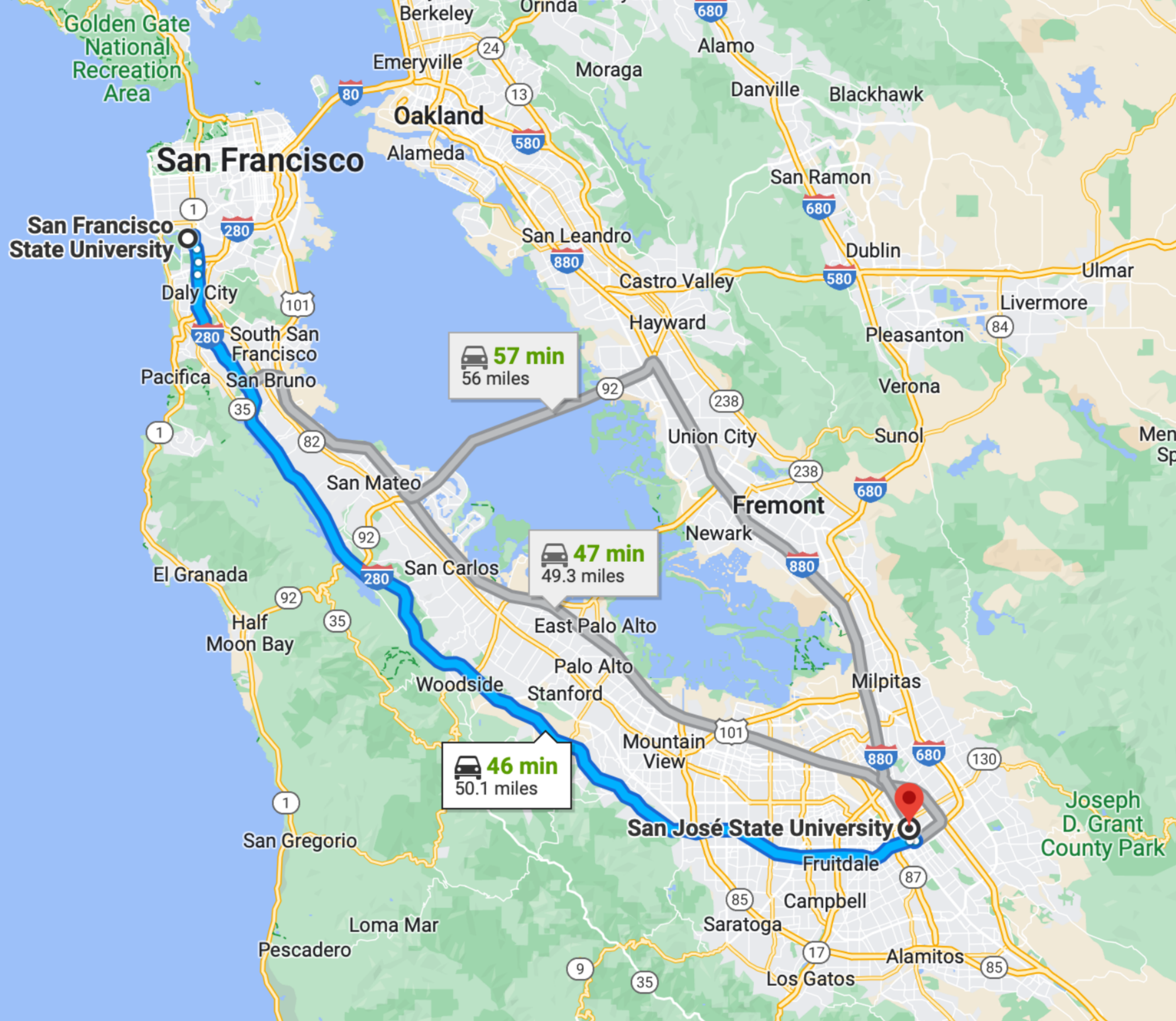
**Corollary:** UCS computes the min-cost path to the goal node.

# Using domain knowledge

- So far: **uninformed search**
  - Algorithms that don't use problem-specific information
  - **Pro:** completely generic — same algorithm works for all search problems
  - **Con:** can't use useful domain knowledge
- Next: **informed search**
  - Use a heuristic function  $h : S \rightarrow \mathbb{R}$  to estimate progress toward goal

# Informed search

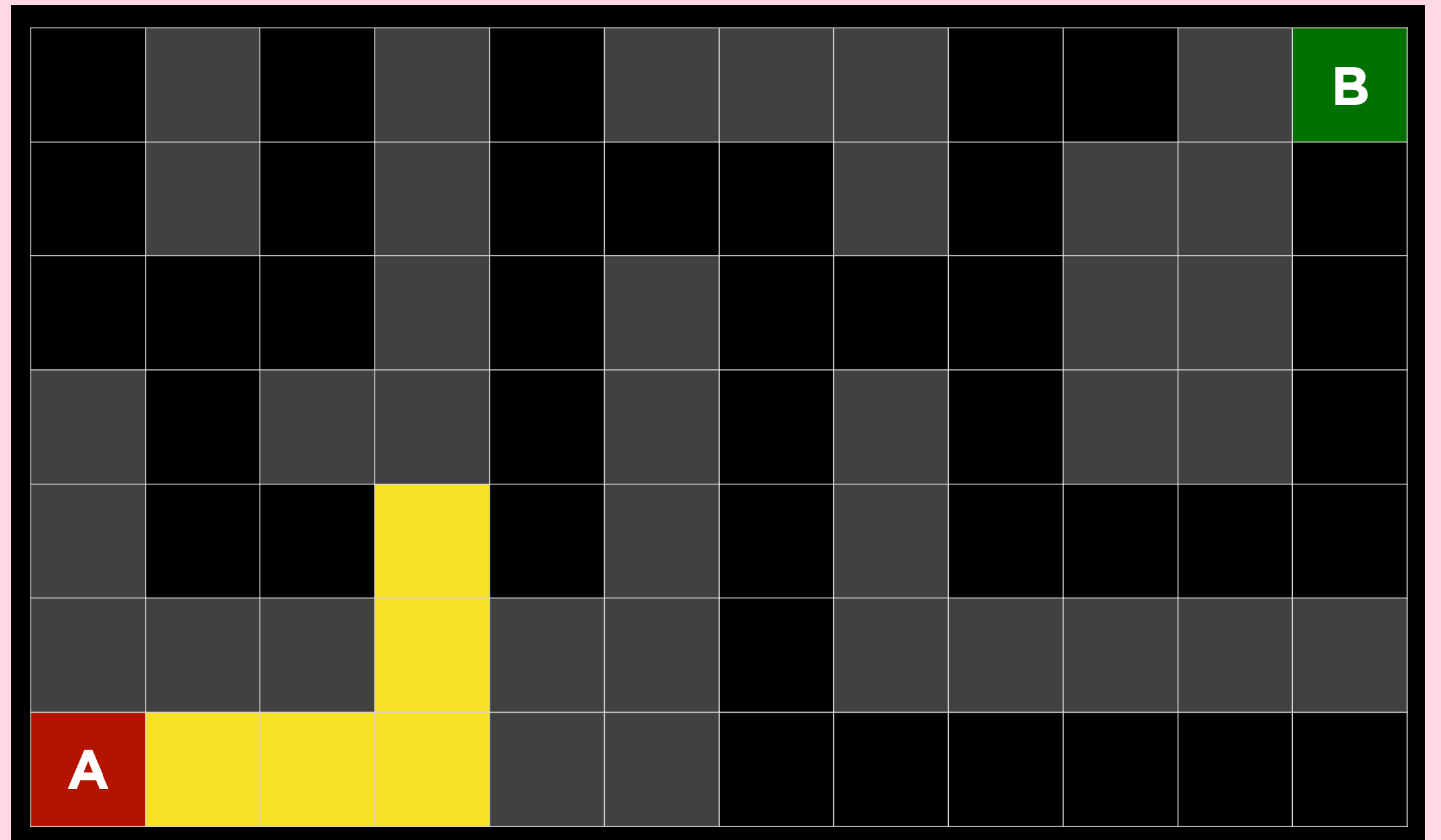


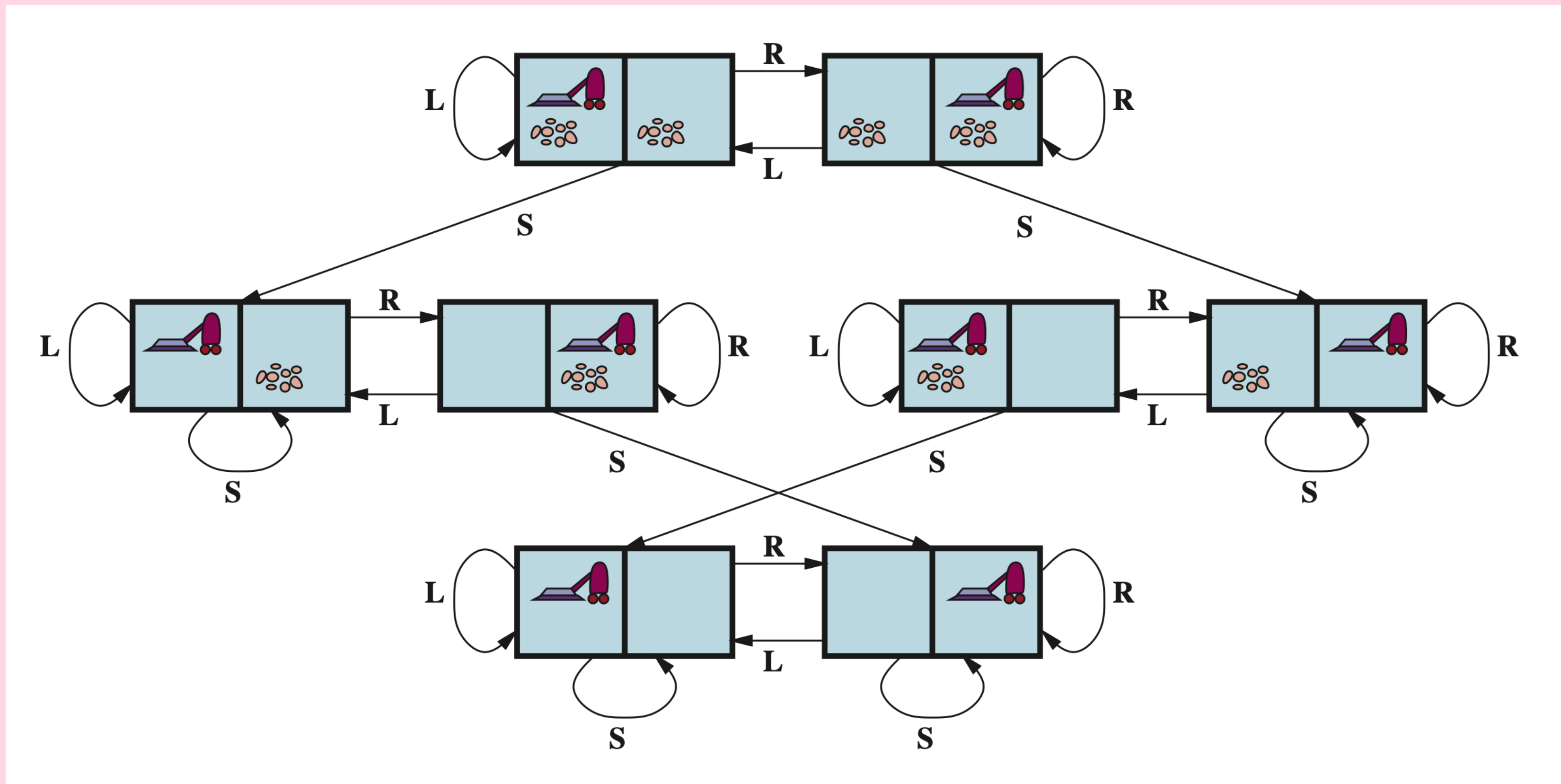


Probably not good  
to start driving  
toward Marin



Probably not good  
to turn left





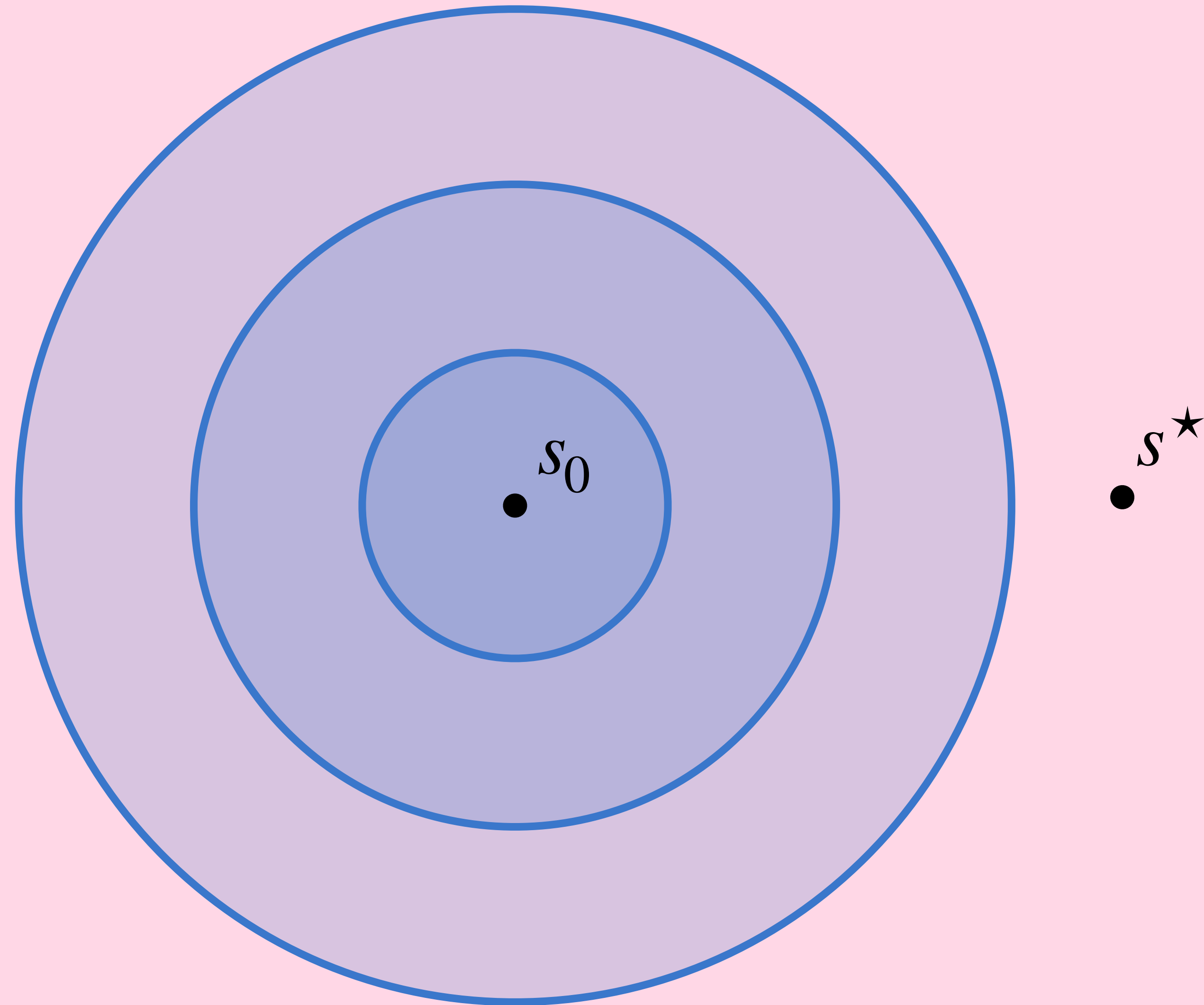
Starting in the upper left state, probably not good to move right before sucking

**How do we know?**



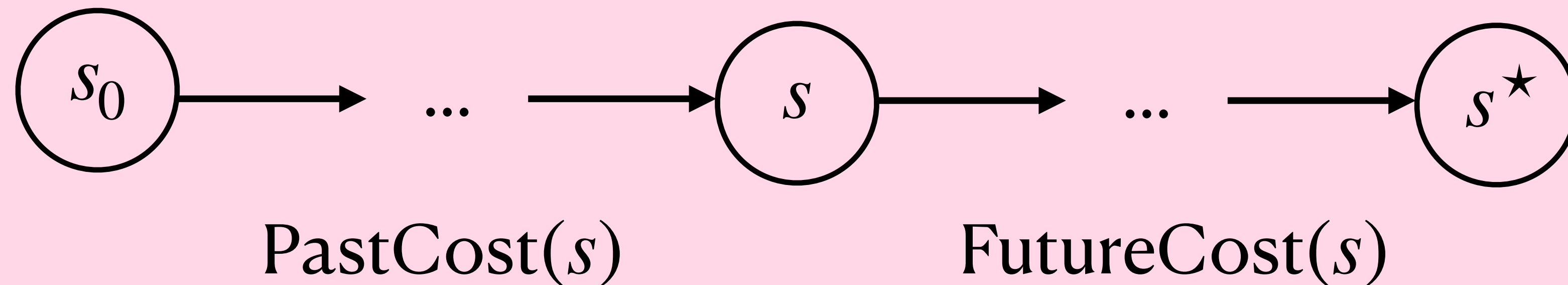
Uniform cost frontier  
is a good idea.

But why bother  
searching in this  
direction?



# Heuristic functions

Consider getting from  $s_0$  to  $s^\star$  on a path through  $s$ .



UCS and BFS work by maintaining a frontier of **uniform PastCost**.

**FutureCost is unknown**, otherwise we could immediately find an optimal solution.

But we can **estimate**  $\text{FutureCost}(s)$  with a simple **heuristic**  $h(s)$ .

# Naïve Idea

- If we had access to FutureCost, then an optimal algorithm is to always expand the node that minimizes FutureCost.
- If all we have is an estimate  $h$ , then why not pick the node that minimizes  $h$ ?
- This is called **greedy search**.

***[greedy search examples]***

# A\* Search

## UCS

- Maintains a frontier of uniform PastCost
- **Correct** but **slow**.

## Greedy search

- Chooses the node that minimizes  $h$
- **Incorrect** but **potentially fast**.

## A\*

- Maintains a frontier of uniform PastCost +  $h$
- **Sometimes correct** and **potentially fast**.

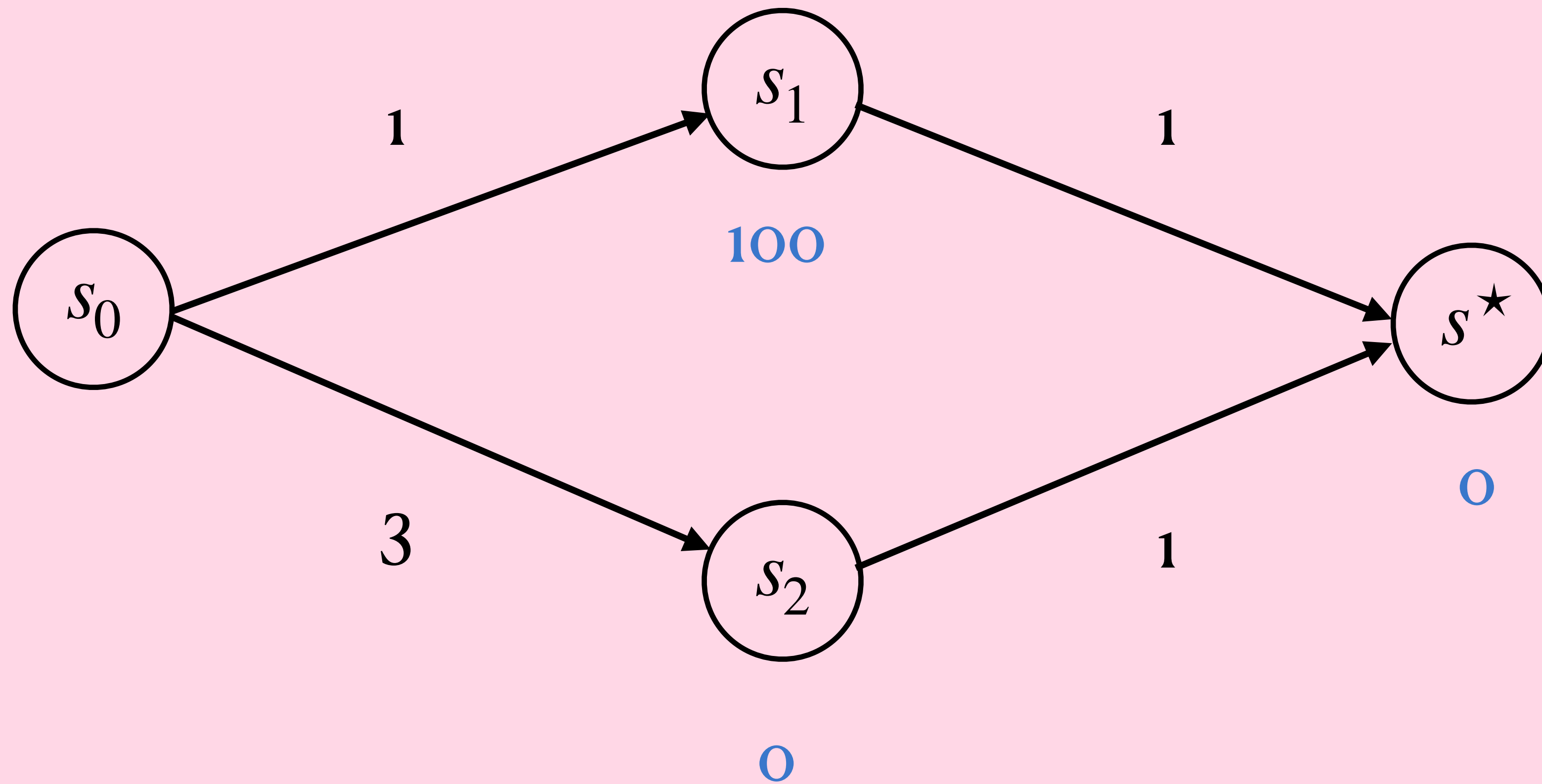
# A\* vs. Greedy

**Problem:** short-term greediness can get you into long-term trouble (true for all greedy algorithms in computer science and in life).

**Key insight:** computing PastCost is easy (just accumulate edge costs), and helps us realize when a prior greedy decision has led us astray.

***[A\* maze example]***

# $A^*$ can be wrong



Action costs

$h$